

# Optimal Peer Selection in a Free Market Peer-Resource Economy

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# Introduction

- Problem of free-riders in P2P Systems.
- Need for a compelling incentive system.

# Free Market Peer Resource Economy

- Online Market Place - Peers buy and sell surplus resources.
- Sellers free to set prices and Buyers "shop" for the best deals.
- Example ?
- Focus: File sharing systems
  - Downloading & Streaming Apps.

# Outline

- Pricing Model
- Problem of optimal peer selection
- Analysis : P2P Downloading
- Analysis : P2P Streaming

# Pricing model : Goods

Goods which can be sold :

- Content
- Resources for content delivery

# Pricing model : Content Delivery Resources

Factors effecting seller's price :

- Rate of transfer
- Duration of transfer

## Pricing model (3)

- We consider pricing functions of the form

$$\text{price} = c_i(b_i) t_i$$

where  $b_i$  is the rate and  $t_i$  transfer duration

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- Analysis : P2P Streaming

# Problem of Optimal Peer Selection

- How to select a subset of sellers so that some optimization criteria are satisfied ?
- Typical optimization criteria could be ?
  - Minimize cost
  - Minimize latency
  - Maximize service reliability.

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## P2P Downloading: Customer's Perspective

- ❑ Customer's objective is to download the file in minimum time possible.
- ❑ Server's price increases with rate.
- ❑ Customer has a limited budget

# Download Protocol

- Search in the overlay identifies, say,  $I$  seller peers offering the requested file.
- Each seller peer advertises its pricing function  $c_i(b)$  and range of offered download rates  $[l_i, u_i]$
- Customer chooses sellers and then downloads file chunks in parallel from the chosen sellers

# P2P Downloading : Optimization Problem

Minimize:

delay of downloading file

Subject to:

- (i) downloading cost is  $\leq$  customer's budget =K
- (ii) file, F bytes, is completely downloaded

# Review of notation

- $I$  servers have file
- $[l_i, u_i]$  = set of transfer rates available from server  $i$
- $b_i$  = chosen transfer rate from server  $i$
- $t_i$  = duration of chunk transfer from server  $i$
- $c_i(b_i) t_i$  = server  $i$ 's price
- $K$  = buyer's budget
- $F$  = file size

# P2P Downloading : Optimization Problem (2)

$$\min \max \{t_1, t_2, \dots, t_I\}$$

subject to

$$c_1(b_1)t_1 + c_2(b_2)t_2 + \dots + c_I(b_I)t_I \leq K$$

$$b_1t_1 + b_2t_2 + \dots + b_I t_I = F$$

$$l_i \leq b_i \leq u_i \quad i = \{1, 2, \dots, I\}$$

$$t_i \geq 0 \quad i = \{1, 2, \dots, I\}$$

# P2P Downloading : Lemma

**Lemma 1** : Suppose  $c_i(b)/b$  is non-increasing.  
There exists an optimal solution with the property  $b_i = u_i$  for all indices  $i$  such that  $t_i > 0$ .

This lemma implies that we may set  $b_i = u_i$  for all  $i$ , thus

w.l.o.g                       $c_i := c_i(u_i)/u_i$

# P2P Downloading : Optimization Problem (2)

$$\min \max \{t_1, t_2, \dots, t_I\}$$

subject to

$$c_1 t_1 + c_2 t_2 + \dots + c_I t_I \leq K$$

$$u_1 t_1 + u_2 t_2 + \dots + u_I t_I = F$$

$$t_i \geq 0 \quad i = \{1, 2, \dots, I\}$$

# P2P Downloading : Solution(1)

Define  $B_j = \sum_{k=1}^j u_k$  and  $\beta_j = \sum_{k=1}^j c_i u_k$

□ Case 1 : If customer peer has large enough budget

$\Rightarrow \frac{K}{F} \geq \frac{\beta_I}{B_I}$  then buy all bandwidth offered

Solution:

$$t_i = \frac{F}{\sum_{k=1}^I u_k}$$

# P2P Downloading : Solution(1)

Define  $B_j = \sum_{k=1}^j u_k$  and  $\beta_j = \sum_{k=1}^j c_k u_k$

□ Case 2 : If customer peer has a smaller budget

$$\frac{\beta_j}{B_j} \geq \frac{K}{F} \geq \frac{\beta_{j-1}}{B_{j-1}}, j \in [1, I]$$

Solution:

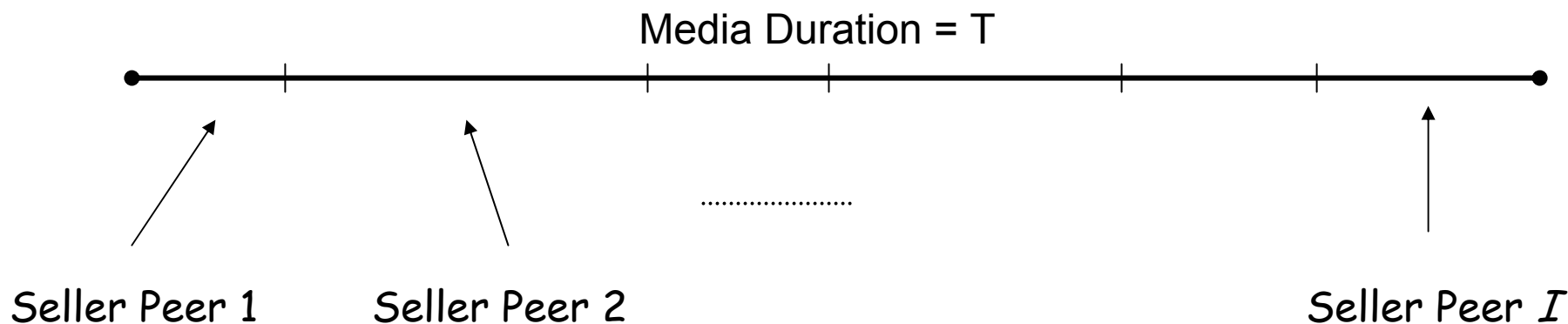
$$t_i = \begin{cases} \frac{Fc_j - K}{c_j B_j - \beta_j} & i \leq j-1 \\ F - B_{j-1} t_{j-1} & i = j \\ 0 & i > j \end{cases}$$

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# P2P Streaming: Possible Approaches(1)

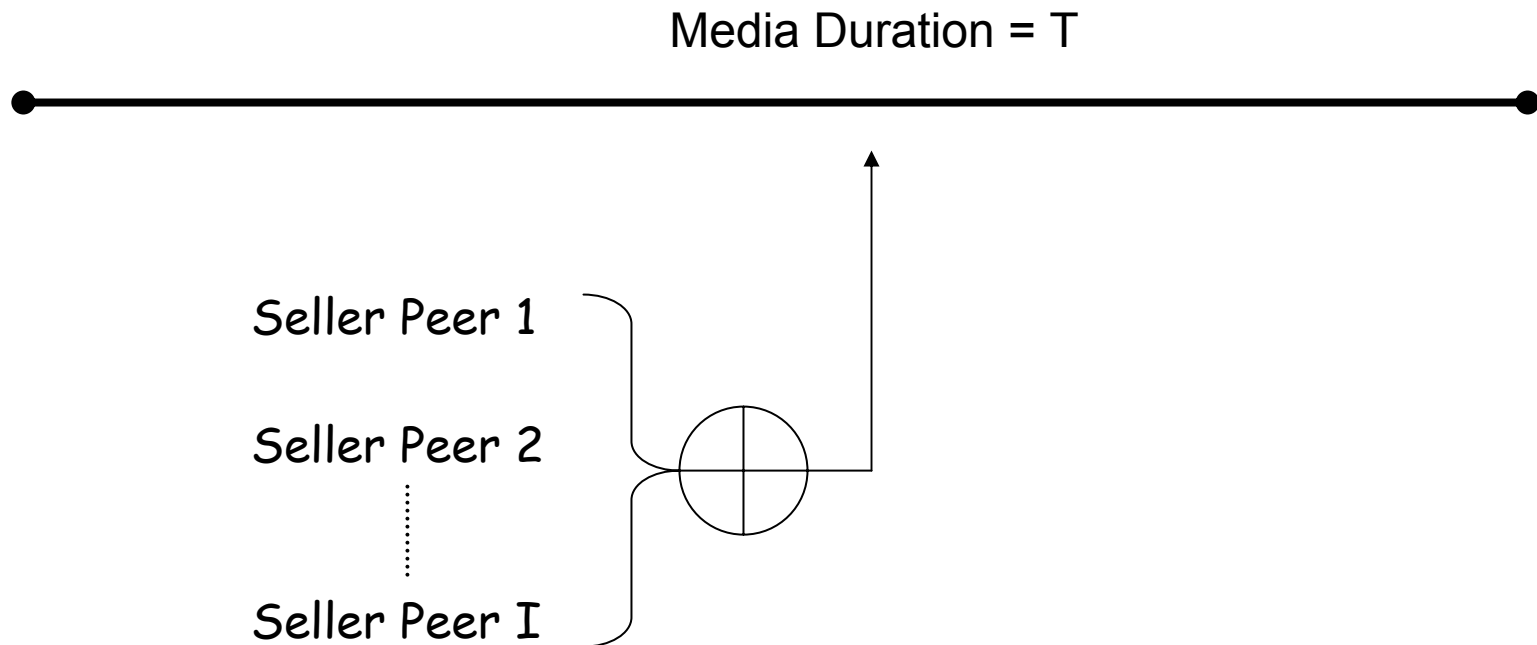
- Time Segmentation
  - Video partitioned along the time axis.
  - Seller peers stream one segment each.
  - Customer peer buffers bytes until scheduled play-out times.



# P2P Streaming: Possible Approaches(2)

## □ Rate Segmentation

- Each seller peer contributes bytes from each frame



# P2P Streaming: Rate vs. Time Segmentation

**Theorem** : If  $c_i(b)$  grow at least linearly, then for any solution  $S$  that uses time segmentation, there is a solution  $S'$  using rate segmentation that has no larger cost.

# P2P Streaming : Optimization Problem

**Minimize:**

**cost of streaming**

**Subject to:**

- (i) **continuous playback at rate  $r$**
- (ii) **even if one server peer goes down**

# P2P Streaming : Optimization Problem(2)

$$\min [c_1(b_1) + c_2(b_2) + \dots + c_I(b_I)]^* T$$

s.t.

$$\underbrace{\sum_{\substack{i=1 \\ i \neq j}}^I b_i}_{i \neq j} \geq r, j = 1, 2, \dots, I$$

# P2P Streaming: Optimization Problem (3)

Alternatively

$$\min C(y), \quad y \in (0, r)$$

Where,

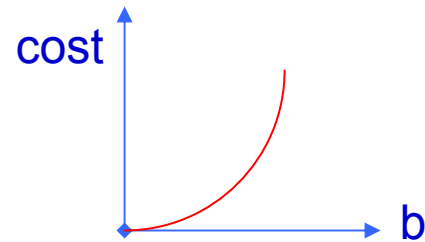
$$C(y) = \min [c_1(b_1) + c_2(b_2) + \dots + c_I(b_I)]$$

$$b_1 + b_2 + \dots + b_I \geq r + y$$

$$b_i \leq y, \quad i = 1, 2, \dots, I$$

# P2P Streaming: convex costs

- **Theorem:** If all  $c_i(b)$  are convex then  $C(y)$  is also a convex function.

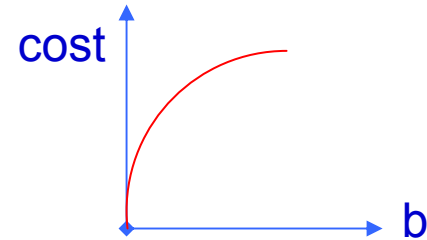


- **Solution Outline:**

- For a given  $y$ ,  $C(y)$  is generated through marginal allocation.
- Streaming problem is greedily solved by finding  $\min C(y)$ ,  $0 \leq y \leq r$

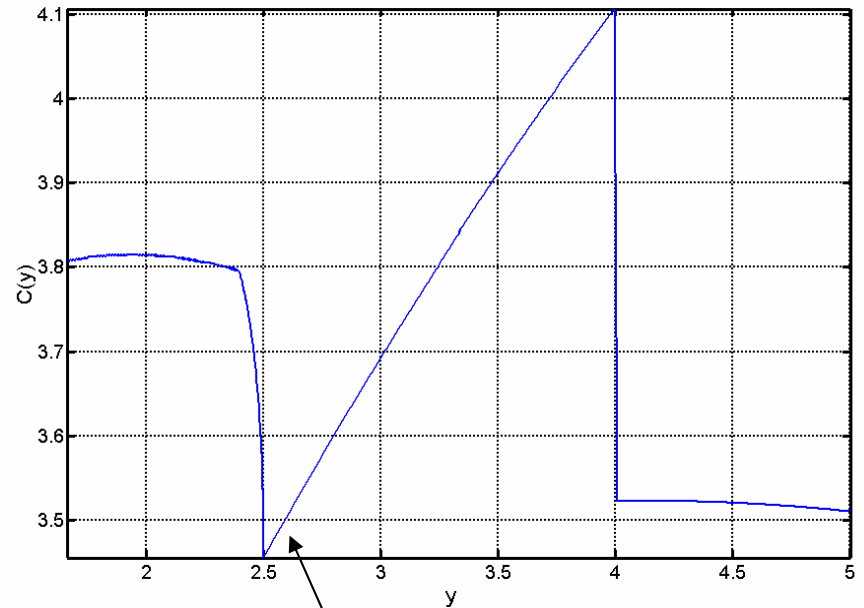
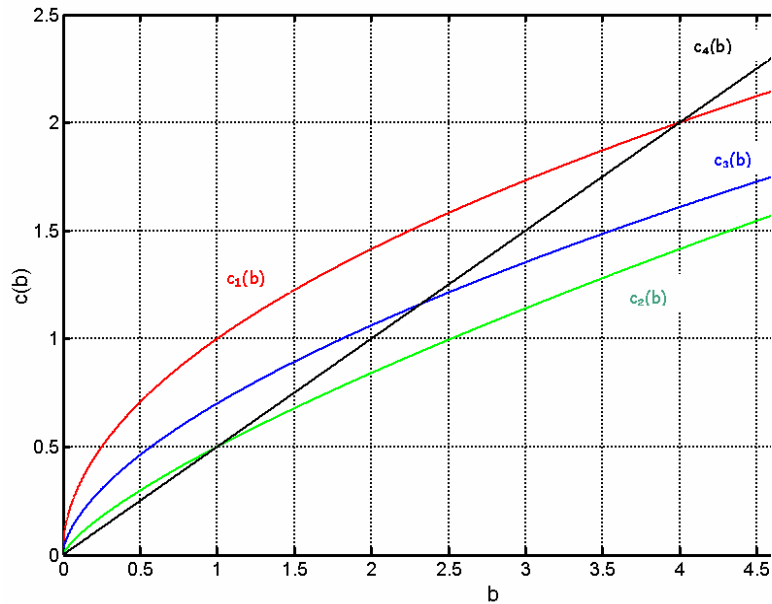
# P2P Streaming: concave costs

- **Theorem:** If all  $c_i(b)$  concave, then we only need to consider the extreme points of feasible region.



- **Solution Outline:**
  - For a given  $y$ , solve  $C(y)$  by computing minimum objective value at the extreme points.
  - Streaming problem is solved by line search to  $\min C(y)$ ,  $0 \leq y \leq r$

# P2P Streaming: concave costs - example



$$c_1(b) = \sqrt{b}$$

$$c_2(b) = 0.5b^{\frac{3}{4}}$$

$$c_3(b) = 0.7b^{\frac{3}{5}}$$

$$c_4(b) = 0.5b$$

$$r = 5$$

$\min C(y) = C(2.5)$

# Summary

- Use of free market resource economy model to provide incentives.
- Customer can make purchase decisions through an Agent
- Discussion for two applications : downloading and streaming.

Questions and Comments ?